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TITLE: Decoding method for trellis codes with large free distances

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## INVENTOR-INFORMATION:

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US-CL-CURRENT: 375/341; 375/265, 379/43

## CLAIMS:

What is claimed is:

1. A decoding method for a trellis code  $T$  of which the encoding can be implemented by first using an encoder of a binary trellis code  $C_{sup}(1)$  to encode an  $r$ -bit message  $u(t)$  into an intermediate binary  $m$ -tuple  $v_{sup}(1)(t)$  at a  $t$ -th time unit which is sequentially converted into binary  $m$ -tuples  $v_{sup}(2)(t), \dots, v_{sup}(L)(t)$  and a signal point  $v_{sup}(L+1)(t)$  in a signal space  $\Omega$  through  $L$  ( $L \geq 2$ ) processors, comprising the decoding steps of:

(a) at a  $(t + \text{LAMBDA}_{sup}(1) + \dots + \text{LAMBDA}_{sup}(L))$ -th time unit, a processor  $P_{sup}(L)$  determining a set  $T_{sub.M.sup}(L)(t + \text{LAMBDA}_{sup}(1) + \dots + \text{LAMBDA}_{sup}(L-1))$  which consists of  $2_{sup.m}$  metrics for the  $2_{sup.m}$  possible values of  $v_{sup}(L)(t + \text{LAMBDA}_{sup}(1) + \dots + \text{LAMBDA}_{sup}(L-1))$ , based on part of the set  $\{y(t+i): i \text{ ltoreq } \text{LAMBDA}_{sup}(1) + \dots + \text{LAMBDA}_{sup}(L)\}$  and part of the set  $\{v_{sup}(L)(t-j): j \text{ gtoreq } \text{lambda}\}$ ; feeding the metric set  $T_{sub.M.sup}(L)(t + \text{LAMBDA}_{sup}(1) + \dots + \text{LAMBDA}_{sup}(L-1))$  into a processor  $P_{sup}(L-1)$ ; wherein  $y(t+i)$  is a received symbol to be decoded,  $\text{LAMBDA}_{sup}(1), \dots, \text{LAMBDA}_{sup}(L)$  are nonnegative constants and  $\text{lambda}$  is a truncation length to be used in decoding  $C_{sup}(1)$ ;

(b) for  $l=L-1, L-2, \dots, 1$ , a processor  $P_{sup}(l)$  determining a set  $T_{sub.M.sup}(l)(t + \text{LAMBDA}_{sup}(1) + \dots + \text{LAMBDA}_{sup}(l-1))$  which consists of  $2_{sup.m}$  metrics for the  $2_{sup.m}$  possible values of  $v_{sup}(l)(t + \text{LAMBDA}_{sup}(1) + \dots + \text{LAMBDA}_{sup}(l-1))$  based on part of  $\{T_{sub.M.sup}(l+1)(t+i): i \text{ ltoreq } \text{LAMBDA}_{sup}(1) + \dots + \text{LAMBDA}_{sup}(l)\}$  and part of the set  $\{v_{sup}(l)(t-j): j \text{ gtoreq } \text{lambda}\}$ ; feeding the metric set  $T_{sub.M.sup}(l)(t + \text{LAMBDA}_{sup}(1) + \dots + \text{LAMBDA}_{sup}(l-1))$  into a processor  $P_{sup}(l-1)$ ; wherein  $\text{LAMBDA}_{sup}(1)$  is a nonnegative constant;

(c) a processor  $P_{sup}(0)$  recovering the transmitted symbol  $u(t - \text{lambda} + 1)$  and  $v_{sup}(1)(t - \text{lambda} + 1)$  by applying the Viterbi decoding algorithm to the trellis of  $C_{sup}(1)$  and using part of the set  $\{T_{sub.M.sup}(1)(t-i): i \text{ gtoreq } 0\}$ ; feeding back the recovered  $v_{sup}(1)(t - \text{lambda} + 1)$  to  $P_{sup}(1)$ ;

(d) for  $l=1, 2, \dots, L-1$ ,  $P_{sup}(l)$  processing the set  $\{v_{sup}(l)(t-j): j \text{ gtoreq } \text{lambda} - 1\}$  and determining  $v_{sup}(l+1)(t - \text{lambda} + 1)$  which is then fed back to  $P_{sup}(l+1)$ .

2. A decoding method as in claim 1, wherein the encoder of said single binary trellis

code  $C_{sup}(1)$  is replaced by the encoders of a plurality of binary trellis codes which together convert  $u(t)$  into  $v_{sup}(1)(t) = (v_{sub.1.sup}(1)(t), v_{sub.2.sup}(1)(t), \dots, v_{sub.m.sup}(1)(t))$ , which is then processed by  $L$  multilevel delay processors and  $L$  signal mappers; and wherein said Viterbi decoding algorithm for  $C_{sup}(1)$  used in the processor  $P_{sup}(0)$  is replaced by a plurality of Viterbi decoding algorithms for said plurality of binary trellis codes.

3. A decoding method as in claims 1 or 2, wherein the encoding of said trellis code  $T$  can be implemented by the steps of:

i) using the encoder of a single binary trellis code  $C_{sup}(1)$  to encode an  $r$ -bit message  $u(t)$  at the  $t$ -th time unit so as to obtain an  $m$ -bit output  $v_{sup}(1)(t) = (v_{sub.1.sup}(1)(t), v_{sub.2.sup}(1)(t), \dots, v_{sub.m.sup}(1)(t))$ ;

ii) processing said output of the encoder through a multilevel delay processor  $Q_{sup}(1)$  to obtain a processed output  $s_{sup}(1)(t) = (s_{sub.1.sup}(1)(t), s_{sub.2.sup}(1)(t), \dots, s_{sub.m.sup}(1)(t))$  such that ##EQU11## wherein  $\lambda_{sub.1.sup}(1), \lambda_{sub.2.sup}(1), \dots, \lambda_{sub.m.sup}(1)$  are nonnegative constants;

iii) feeding said output of  $Q_{sup}(1)$  into a signal mapper  $S_{sup}(1)$  to select a signal point  $w_{sup}(1)(s_{sup}(1)(t)) = v_{sup}(2)(t) = (v_{sub.1.sup}(2)(t), v_{sub.2.sup}(2)(t), \dots, v_{sub.m.sup}(2)(t))$ ;

iv) for  $l=2, 3, \dots, L$ , feeding the output  $v_{sup}(1)(t)$  into a multilevel delay processor  $Q_{sup}(l)$  to obtain a processed output  $s_{sup}(l)(t) = (s_{sub.1.sup}(l)(t), s_{sub.2.sup}(l)(t), \dots, s_{sub.m.sup}(l)(t))$ , such that ##EQU12## wherein  $\lambda_{sub.1.sup}(l), \lambda_{sub.2.sup}(l), \dots, \lambda_{sub.m.sup}(l)$  are nonnegative constants, and then feeding  $s_{sup}(l)(t)$  into a signal mapper  $S_{sup}(l)$  to select a signal point  $w_{sup}(l)(s_{sup}(l)(t)) = v_{sup}(l+1)(t)$ , whereby if  $l=L$ . A decoding method as in claim 3, wherein the parameter  $\lambda_{sup}(l)$  of the processor  $P_{sup}(l)$  is no less than  $\Sigma_{i=1}^m \lambda_{sub.i.sup}(l)$  for  $1 \leq l \leq L$ .

5. A decoding method as in claim 3, wherein said constants  $\lambda_{sub.1.sup}(1), \lambda_{sub.2.sup}(1), \dots, \lambda_{sub.m.sup}(1)$  are distinct values and/or at least one of the constants  $\lambda_{sub.1.sup}(1), \lambda_{sub.2.sup}(1), \dots, \lambda_{sub.m.sup}(1)$  is not zero.

6. A decoding method as in claim 3, wherein said constants  $\lambda_{sub.1.sup}(1), \lambda_{sub.2.sup}(1), \dots, \lambda_{sub.m-1.sup}(1)$  are all identical to a constant  $\lambda_{sup}(1)$  and  $\lambda_{sub.m.sup}(1) = 0$ .

7. A decoding method as in claims 1 or 2, wherein said binary trellis code or a plurality of binary trellis codes is/are a linear and time-invariant code or a plurality of linear and time-invariant codes respectively, i.e., a binary convolutional code or a plurality of binary convolutional codes.

8. A decoding method as in claims 1 or 2, wherein the signal space  $\Omega$  is a signal constellation; then said method can be used to decode trellis coded modulation.

9. A decoding method as in claims 1 or 2, wherein the signal space  $\Omega$  is a collection of binary  $m$ -tuples; then said method can be used to decode binary trellis codes.

10. A decoding method as in claims 1 or 2, wherein the  $r$ -bit information,  $u(t)$ , represents  $m/m_{sup}(1)$  binary  $(r \cdot \text{multidot} \cdot m_{sup}(1) / m)$ -tuples, whereby  $m_{sup}(1)$  is a positive integer.

11. A decoding method as in claims 1 or 2, wherein the binary  $m$ -tuple  $v_{sup}(1)(t), 1 \leq t \leq L$ , can be decomposed into  $m/m_{sup}(1)$  binary  $m_{sup}(1)$ -tuples, wherein  $m_{sup}(1)$  is a positive integer; and  $v_{sup}(L+1)(t)$  represents  $m/m_{sup}(L+1)$  signal points in  $\Omega$  which consists of  $2^{(m \cdot \text{sp} \cdot (L+1) \cdot \text{sup})}$  signal points, wherein  $m_{sup}(L+1)$  is a positive integer.